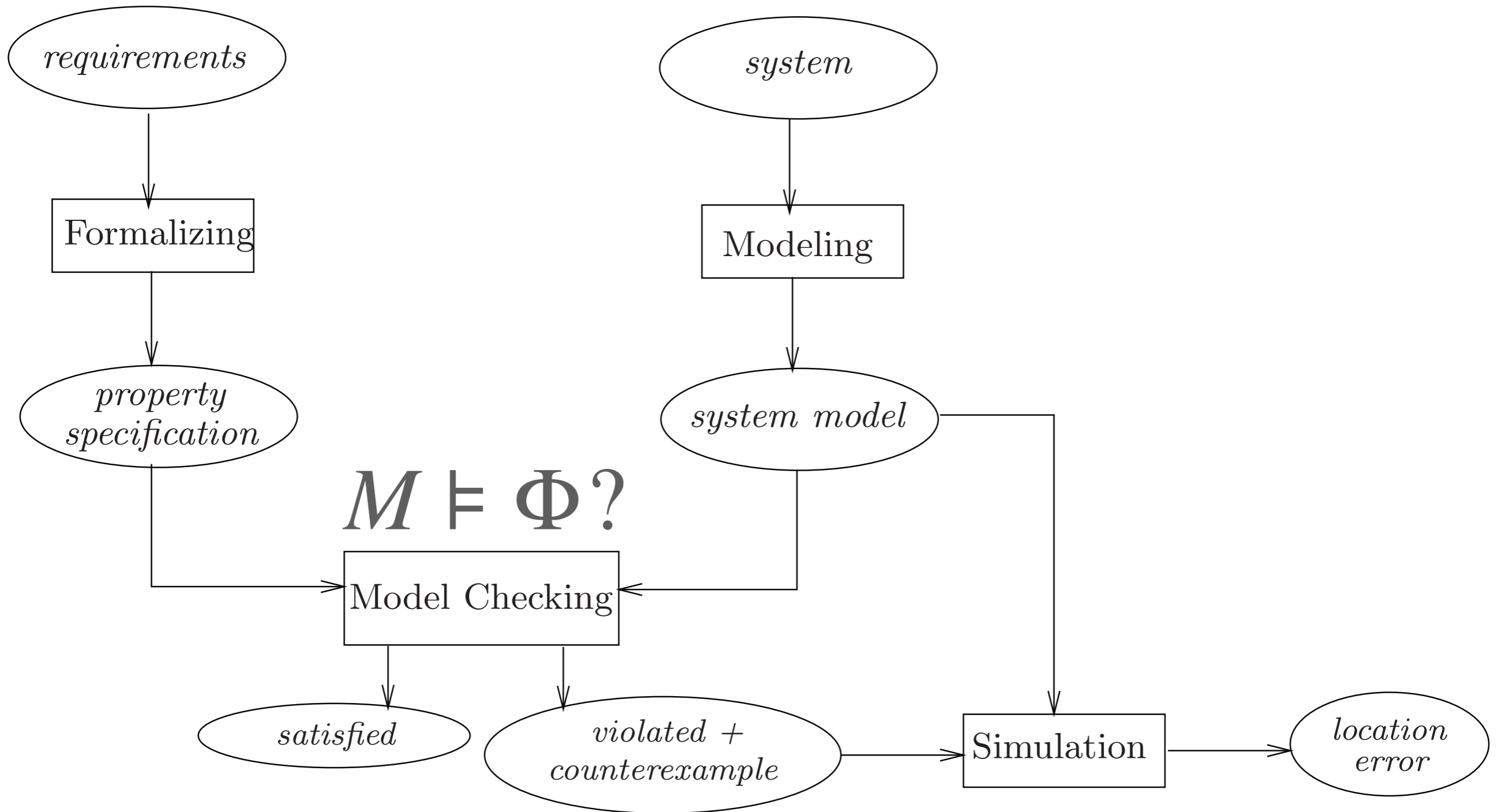


# Model Checking

$$M \models \Phi?$$

## Introduction

# What is Model Checking?



# Model Checking process

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  - formalise the property to be checked using the property specification language.

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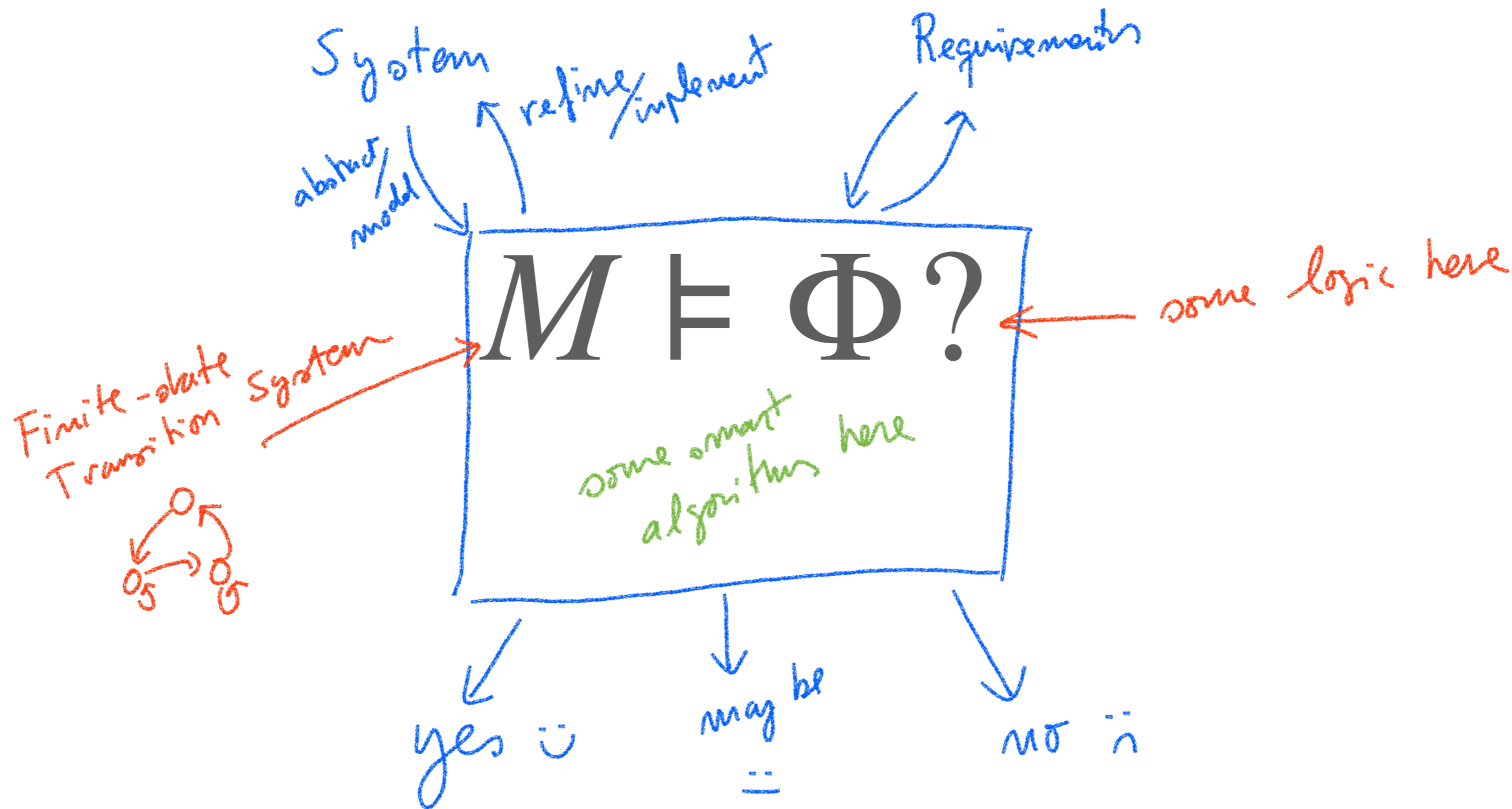
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- *Running* phase:
  - run the model checker to check the validity of the property in the system model.
- *Analysis* phase:
  - property satisfied? → check next property (if any);
  - property violated? →
    1. analyse generated counterexample by simulation;
    2. refine the model, design or property;
    3. repeat the entire procedure.
  - out of memory? → try to reduce the model and try again.



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- *Easy integration* in existing development cycles: its learning curve is not very steep, and empirical studies indicate that it may lead to shorter development times;
- *Sound and mathematical underpinning*, based on theory of graph algorithms, data structures and logic.

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- Correct results not guaranteed: as any other tool, it may contain *software defects*;